

Aspects of CONNES EMBEDDING CONJECTURE (PROBLEM)

Discrete groups (countable)

Let Γ be a discrete group

pb. Given $\varepsilon > 0$, $F \subseteq \Gamma$ finite

find N, Φ (depending on ε, F)

$$\Phi: F \longrightarrow \mathcal{U}(N)$$

such that

- HYPERLINEAR GROUP
- 1) $\|\Phi(\delta_1 \delta_2) - \Phi(\delta_1) \Phi(\delta_2)\|_2 < \varepsilon$
if $\delta_1, \delta_2, \delta_1 \delta_2 \in F$
 - 2) $\|\Phi(\delta) - Id\|_2 > 1 - \varepsilon$ if $\delta \in F, \delta \neq e$

Explanation of terms $\mathcal{U}(N) \subseteq M_N(\mathbb{C})$

is the group of unitary matrices

$$\|x\|_2^2 = \text{tr}(x^*x) \quad (\text{making } M_N(\mathbb{C}) \text{ a Hilbert space})$$

$$\text{tr}(A) = \frac{1}{N} \text{Tr}(A) = \frac{1}{N} \sum_{i=1}^N a_{ii}$$

C E. Conjecture : All discrete group
are hyperlinear

stronger form. Replacing $U(N)$ by S_N

S_N group of permutations $\subseteq U(N)$

We get the definition of sofic groups.

(distance becomes Hamming distance)

Original form of definition of Hyperlinear

Assume for simplicity Γ is finitely generated
by g_1, g_2, \dots, g_n . Then we have a length

function $|g| = \inf \{s \mid g = g_{i_1} g_{i_2} \dots g_{i_s}\}$

If $w = g_{i_1} g_{i_2} \dots g_{i_p}$ is any word in

the generators, then for any N

we get a word map (denoted still by w)

$$w : U(N)^n \rightarrow U(N)$$

$$w(U_1, U_2, \dots, U_n) = U_{i_1} U_{i_2} \dots U_{i_p}$$

3) Word maps are object of an intensive study. Questions: When is w surjective? What is \bar{w}^* (Haar measure)? (see Peter Mager) it is Haar measure if $\Gamma = F_n$, w a primitive word)

Reformulation of C. E. Conj.

Given Γ , with generators $\delta_1, \delta_2, \dots, \delta_u$

find, for every $\varepsilon > 0$, every $S_0 \in \mathbb{N}$ (instead of \mathbb{N})

N and U_1, U_2, \dots, U_u such that

for all $w = \delta_{i_1} \dots \delta_{i_s}$, $s \leq S_0$

$$\text{tr}(w(U_1, \dots, U_u)) \approx \begin{cases} \frac{2}{\varepsilon} & 0 \text{ if } w \neq e \\ \frac{2}{\varepsilon} & 1 \text{ if } w = e \end{cases}$$

(The equivalence with previous comes from

$$\|1 - u\|_2^2 = 1 - 2\text{tr}(u) + 1 = 2(1 - \text{tr}(u))$$

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Linearization of Γ .

$$\mathbb{C}(\Gamma) = \left\{ \sum_{g \in \Gamma} a_g g \mid a_g \in \mathbb{C} \right\}$$

all 0 with

exception a finite number

$\mathbb{C}(\Gamma)$ is an algebra: with $g_1 g_2 = (g_1 g_2)$

(product in the algebra = product in the group)

(Ex. If Γ_2 is the free group with generators a, b
then $(a + b + a^{-1} + b^{-1})^2 = \sum_{|w|=2} w + 4$)

$\mathbb{C}(\Gamma)$ has a trace $\bar{\tau}$ (like trace on matrices)

$$\bar{\tau} \left(\sum_i a_g g \right) = a_e$$

It is a trace i.e. $\bar{\tau}(xy) = \bar{\tau}(yx)$ ($\forall x, y \in \mathbb{C}(\Gamma)$)

$$\text{since } \bar{\tau}(g_1 g_2) \neq 0 \Leftrightarrow g_1 g_2 = e \Leftrightarrow g_2 g_1 = e \Leftrightarrow \bar{\tau}(g_1) = 1$$

$\bar{\tau}$ is also positive, i.e. $\bar{\tau}(x^* x) \geq 0$

(Since representing $\mathbb{C}(\Gamma) \subseteq B(\ell^2(\Gamma))$)

by left convolutions

$$\bar{\tau}(x) = \langle x \delta_e, \delta_e \rangle \text{ positive functional}$$

In plain words Gromov Embedding conj
 says: Approximate the trace on the
 group algebra, on finite sets, by
 the traces on matrices.

More precisely, General formulation
 of Gromov Embedding Problem (counterexample
 exists)

Given OT algebra with involution $*$
 (like $C(\Gamma)$ with $(\sum a_g g)^* = \sum \bar{a}_g g^{-1}$)

and trace τ (+ boundedness condition
 ($\| \cdot \|_1$ von Neumann algeb))

$$(\forall) x, (F) M \tau(x^* y^* y x) \leq M \tau(y^* y) \quad (\forall) y \in \mathcal{O}$$

Given $x_1, x_2, \dots, x_n \in \mathcal{O} \quad (\forall) \epsilon_0, \epsilon > 0$

(F) $N \in \mathbb{N}, X_1, X_2, \dots, X_n \in M_N(\mathbb{C}) \leq 1$

$$|\tau(x_{i_1} x_{i_2} \dots x_{i_s}) - \text{tr}(X_{i_1} \dots X_{i_s})| < \epsilon \quad \text{for } s \leq s_0$$

Examples: Case $\Gamma = \langle a, b \rangle$

If Γ is residually finite (i.e. exists

$\Gamma_n \triangleleft \Gamma$ (normal of finite index such that

$\Gamma \cap \Gamma_n = e$) then take $\Phi_n: \Gamma \rightarrow B(\ell^2(\Gamma/\Gamma_n))$

the representation by left convolution.

They do the job because $(\forall) g \in \Gamma$ there

exists a quotient where $\Phi_n(g)$ acts by

permutation with no fixed elements, so it

has zero trace.

Ex $\Gamma = F_{\mathbb{R}} = \langle a_1, \dots, a_n \rangle$ (it is residually finite)

Th. (Voiculescu). If $X_i = \text{Re } a_i = \frac{1}{2}(a_i + a_i^*)$

Then random matrices ^{size N} with independent gaussian entries have the property

that expected values of traces of monomials in matrices converge to the trace τ of corresponding monomials in the algebra

(7) Voiculescu's
 Free entropy = normalized volume (as if...)
 of microstates (the matrices X_1, X_2, \dots, X_n)
 $\chi(x_1, \dots, x_n) = n$ for x_1, x_2, \dots, x_n from \overline{F}_n

For C.E. Conjecture: Possible strategy

Take $w_1, w_2 \in \overline{F}_n$ such that w_2
 is not a consequence of w_1 (i.e.

$w_2 \neq$ product of conjugates of w_1)

Obs Then C.E. Conj \Leftrightarrow $(\exists \varepsilon > 0)$ $(\exists) U_1, U_2, \dots, U_n \in U(\mathbb{H})$
 unitaries such that

$$|\operatorname{tr}(w_1(U_1 \dots U_n))| > 1 - \varepsilon$$

$$|\operatorname{tr}(w_2(U_1 \dots U_n))| < \varepsilon$$

This strategy would lead to

\overline{F}_n / w_1 is hyperlinear
 (relate)

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Th (R). The moments

$$\int \text{tr}(\omega_1)^\alpha \text{tr}(\omega_2)^\beta d\mu_{H_{\alpha+\beta}} \text{ are}$$

precisely computable if $N > \alpha|\omega_1| + \beta|\omega_2|$

Not sufficient for solving since $\text{tr}(\omega)$ are

very "small" : $\lim_{N \rightarrow \infty} \frac{\mu(|\text{tr}(\omega(u_1 \dots u_n))| > \epsilon)}{d(N)^n}$

is 0 if ω is untrivial (Th. Voiculescu)

• Hilbert 17th (non-commutative form)

Given $P(x_1, \dots, x_n)$ non-commutative polynomial (assume $x_i = x_i^*$), $P = P^*$

$C \in P \Leftrightarrow \exists H \text{ tr}(P(X_1, X_2, \dots, X_n)) \geq 0$ (A)

(A) X_1, X_2, \dots, X_n matrices $\Rightarrow \exists \mathcal{A} \text{ tr}(P(x_1, \dots, x_n)) \geq 0$ (A)
? \mathcal{A}
(A) \mathcal{A} with trace, $x_i \dots x_n \in \mathcal{A}$

Of course such a p that is of the

$$\text{form } p = \sum_i z_i^* z_i + \sum_j [r_j, a_j]$$

has the property B . (ie: $\bar{\sigma}(p) \geq 0$)

Problem (Hilbert) solved by Artin

Th A positive polynomial is a sum of squares of rationals.

Ex: (commutative variables) (Mokobodzki polynomial)

$$1 + x_1^2 x_2^4 + x_1^4 x_2^2 - 3 x_1^2 x_2^2 =$$

$$= \frac{(x_1^3 x_2 + x_2^3 x_1 - 2 x_1 x_2)^2 (1 + x_1^2 + x_2^2) + (x_1^2 - x_2^2)^2}{(x_1^2 + x_2^2)^2}$$

Th (McCullough Helton)

If $p(x_1, \dots, x_n) \geq 0$

(*) x_1, \dots, x_n matrices of controlled size

$$\Rightarrow p = \sum_i z_i^* z_i$$

(9)

Trace positivity. Tu (Process + others)

If $tr(p(x_1, x_2, \dots, x_n)) \geq 0 \quad (\forall) x_1, x_2, \dots, x_n \in M_n(\mathbb{C})$

then $p = \text{S.O.S} + \text{commutators}$

modulo the fundamental Cayley Hamilton identity in $M_n(\mathbb{C})$

$Tu(\mathbb{R})$. There exists a S.O.S

expression for (B) in terms of

analytic entire non-commutative series

Many better generalizations (Klepp, ...)

Also (Parillo, others) have SDP methods
M. Laurent

for determining $C \geq 0$ st, $p + c$ is a S.O.S

if modulo commutators

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Non-local games formulation (Tsirelson problem)

Heuristically. If two algebras $A, B \subseteq B(H)$ commute: determine how far are these from a situation of the form $A \subseteq B(H_1), B \subseteq B(H_2), B(H) = B(H_1) \otimes B(H_2)$ (tensor product commutation)

Simplified version. A is generated by E_1, E_2, \dots, E_n

B is generated by F_1, F_2, \dots, F_s , where $E_1, \dots, E_n,$

F_1, F_2, \dots, F_s are projections.

Consider \mathcal{C} (convex set of entangled states

$$\mathcal{C} = \left\{ (a_{ij}) = \left\langle \left(\sum_i E_i \otimes F_j \right) \xi, \xi \right\rangle \right\}_{ij} \quad \left. \begin{array}{l} \text{where } A, B \subseteq B(H) \\ \xi \in H, \|\xi\| = 1 \end{array} \right\}$$

How far is this from

$$\mathcal{C}^{\dagger} = \text{same but, } A \subseteq B(H_1), B \subseteq B(H_2) \\ \xi \in H_1 \otimes H_2$$

Non-local games strategy (Hence, =)
proved that describing e^+ would require
a Turing machine that solves all halting
problems (which is impossible), while
 e^{nc} can be determined by SDP methods

First part: J. Zhengeng, A. Vidick,
T. Wright, A. Natarajan, H. Yuen in 2020

C^+ part solved by Navasques, Pinouid

So $C^*(F_n) \otimes C^*(F_n)$ has not a
unique norm so CEP is disproved.

Many open problems remain

1) C.E. Conjecture for discrete groups

2) Give a positivity Criteria for concrete polynomials
e.g coefficients of t in $\sum (e^{x+ty})$

if x, y are positive. This is positive on matrices

(This is BESSIS-MOUSSA-VILLANI tree conjecture
solved by STAHL in 2011)